

Lab 01

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Exercises

Design Thinking

- **Definition.** well-define special keywords like *exhaustive search*, *floor*, *square*.
- **Break.** Break the big problem into subproblems.
- **Partial Progress.** Solve subproblems or relaxed versions.
- **Concrete examples.** Try on concrete cases.
- **Generalize.** Spot the pattern generalizable on any case conforming to the general definition.
- **Connect Ideas.** Figure whether the different ideas and solutions can combined.

Algorithmic Hints

- Are there redundant computations?

1.1.4

Hints

- You are given a square number n . Given some integer k , How can we verify it is the root?
- Follow the exhaustive search strategy, to find the root of n .
- You are given a real number r . Given some integer k , How can we verify it is the floor of r ?
- Follow the exhaustive search strategy, to find the floor of n .
- Combine all previous hints to find a unique definition of $\lfloor \sqrt{n} \rfloor$.
- Follow the exhaustive search strategy, to solve the main problem.

Solution

```
for i in n-1 .. 0
  if (i)^2 <= n
    return i
```

1.1.8

Hints

- Try this case on concrete examples like $m = 2$ and $n = 3$.
- Why $m \bmod n = m$ when $m < n$?

- Recall the definition of $x \bmod n$. What are the possible ranges of $x \bmod n$ for any integer x ?

Solution

It shall swap them as $r = m \bmod n = m$ when $m < n$.

Only once. Given $m > n$, Necessarily $n > m \bmod n$.

1.2.5

Hints

- Convert a concrete decimal number to binary. Observe how the right most digit from the binary representation is obtained.
- Given a binary representation, What is the number we divide on it, so that the quotient eliminate the right most digit?
- Follow the *Decrease and Conquer* strategy, with the above two hints, to solve the problem.

Solution

DecToBin(n):

```
# input: integer n
# output: binary representation as a list

# binary representation
l = [ ]

while n != 0:
    # kth digit from right to left
    b.appendLeft( n % 2 )

    # remove the rightmost digit
    # division output is an integer
    n = n/2
```

Algorithm 1 Convert to the binary representation of a given integer

Require: input is integer n

Ensure: output is a list of binary digits

```
1:  $l \leftarrow [ ]$ 
2: while  $n \neq 0$  do
3:    $l \leftarrow [n \bmod 2] \cup l$ .
4:    $n \leftarrow n/2$ 
5: end while
```

1.2.9

Hint

- Are there duplicated computations?
- Are there pairs tested twice?
- Observe $|a - b| = |b - a|$.
- If we checked all elements with $A[i]$, Do we need to test $A[j]$ with $A[i]$?

Solution

```
MinDistance( A ):
  # input: array of size n
  # output: minimum distance between two distinct elements

  dmin = infinity
  for i in 0 .. n-1:
    for j in i+1 .. n-1:
      dis = | A[i] - A[j] |
      if dis < dmin:
        dmin = dis
```

Algorithm 2 Find the minimum distance between two distinct elements in an array

Require: input is array $A[0..n - 1]$ of numbers

Ensure: output is the minimum distance between any two distinct elements

```
1:  $dmin \leftarrow \infty$ 
2: for  $i = 0$  to  $n - 1$  do
3:   for  $j = i + 1$  to  $n - 1$  do
4:      $dis \leftarrow |A[i] - A[j]|$ 
5:     if  $dis < dmin$  then
6:        $dmin \leftarrow dis$ 
7:     end if
8:   end for
9: end for
10: return  $dmin$ 
```

1.3.1

Hints

- if $A[i] == A[j]$ which index shall be counted? What can we conclude about S ?

Solution

a. Tedious to typeset.

b. No. Observe counting only happens when strictly $i < j$. If $A[i] == A[j]$ then the code counts $A[i]$ not $A[j]$. Therefore $A[i]$ shall succeed $A[j]$. In fact equal cells are reversed in the sorted array.

c. No. It does not modify array A but output is a different array S .

1.3.10

FAILED TO SOLVE.

1.4.2

Hint

- For ascendingly ordered array A , Is it possible for the target value t to exist in $A[i..n - 1]$ given the fact $t > A[i]$?
- Use the above hint to prune the search space.
- Which index of the array you think shall prune the greatest search space.

Solution

For target value t :

- a. Access some element x in the array. If $t \neq x$, We can ignore searching in the right/left side of x .
- b. While linear scanning, Terminate the algorithm earlier once some $A[i] > t$.

1.4.10

Hints

- Is it possible for two strings to be anagrams in case they different lengths?
- Is it possible for two strings to be anagrams if one of them has a character not present in the other?
- You can convert a character to its corresponding ascii number. Use that for a cheaper data strucutre.
- the ascii number corresponds to an index.

Solution

Two strings are *anagrams* if and only if they have the same count of characters.

```
AreStringsAnagrams(A, B):  
    # input two strings
```

```

# output True if anagrams and False otherwise

# if lengths are not the same, then not anagrams
if length(A) != length(B):
    return False

# initialize characters counts to zeros for both strings
A_chCount = B_chCount = [ 0 ] * 26

# Count characters in both strings
for ch in A:
    A_chCount[ int(ch) ] = A_chCount[ int(ch) ] + 1

for ch in B:
    B_chCount[ int(ch) ] = B_chCount[ int(ch) ] + 1

# Anagrams if and only if characters count is exactly the same
return A_chCount == B_chCount

```

Algorithm 3 Detect whether two strings are anagrams

Require: input is two strings

Ensure: output True if anagrams and False otherwise

```

1: if  $|A| \neq |B|$  then
2:   return false
3: end if

4:  $A_{count} \leftarrow B_{count} \leftarrow [ ] \times 26$ 

5: for  $ch \in A$  do
6:    $A_{count}[ch] \leftarrow A_{count}[ch] + 1$ 
7: end for

8: for  $ch \in B$  do
9:    $B_{count}[ch] \leftarrow B_{count}[ch] + 1$ 
10: end for

11: return  $A_{count} == B_{count}$ 

```
